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# **OPTIMIZING AN ALUMINUM EXTRUSION PROCESS**

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### ABSTRACT

Minimizing the amount of scrap generated in an aluminum extrusion process. An optimizing model is constructed in order to select the best cutting patterns of aluminum logs and billets of various sizes and shapes. The model applied to real data obtained from an existing extrusion factory in Kuwait. Results from using the suggested model provided substantial reductions in the amount of scrap generated. Using sound mathematical approaches contribute significantly in reducing waste and savings when compared to the existing non scientific techniques.

Keywords: Minimizing, Cutting Patterns, Aluminum Logs, Mathematical Approaches

## **1. INTRODUCTION**

Aluminum is the third most abundant element in the Earth's crust and it constitutes 7.3 percent by mass. The Aluminum industry contributes significantly to the global economy as well as too many individual economies. The industry employs over a million people worldwide. Aluminum smelting is a capital-intensive, technology-driven industry concentrated in a few relatively dominant companies. Aluminum consumption has enjoyed substantial average growth over the last few decades due to general economic growth and to its substitution of other materials. Kuwait is a member of the Gulf Cooperation Council (GCC); the GCC countries will boost their share of global aluminum output to 15-17% by the end of the decade. There are two types of aluminum industries; the first is the extrusion industry where profiles of different sizes, colors and shapes are produced, while the second is the fabrication industry where various products such as windows, fences and doors are designed from aluminum profiles.

In the aluminum extrusion industry, logs and billets are cut using various stock cutting patterns; the amount of scrap generated is dependent on the cutting method used. The Stock Cutting Problem (SCP) is discussed thoroughly in the literature. One of the first articles was presented by Gilmore and Gomory (1961) where integer programming was utilized for the cutting stock problem. The problem compromised a large number of variables which generally makes the computation infeasible. Gilmore and Gomory (1964) examined the cutting stock problems involving two or more dimensions. Haessler (1971) described a heuristic procedure for scheduling production-rolls of paper through a finishing operation to cut them down to finished roll sizes. The objective was to minimize the cost of trim-loss and that of the reprocessing. Covesrdale and Wharton (1976) presented a heuristic procedure for a nonlinear cutting stock problem; the problem was solved using the pattern enumeration technique.

Sumichrast (1986) addressed a scheduling problem in the woven fiber glass industry as an example of the cutting stock problem with the objective of controlling the wasted production capacity rather than wasted material. A heuristic was developed for the purpose of scheduling the production process. Stadtler (1990) used the column generation method of Gilmore and Gomory (1964) for minimizing the amount of scrap generated from fabricated aluminum made for window frames. Krichagina *et al.* (1998) examined the cutting process of sheets in a paper plant. The main objective was to minimize the long-run average cost of paper waste. In this regard, a two step procedure consisting of linear programming and Brownian control was developed.



Liang et al. (2002) applied an evolutionary algorithm (EP) for cutting stock problems with and without contiguity. Results showed that the EP algorithm is more effective and superior when compared to the genetic algorithm used. Parada et al. (2003) proposed a meta-heuristic approach for solving a non-guillotine stock cutting problem. The approach was a combination of the principles of the constructive and evolutive methods. The results showed an error reduction of around 2%. Hifi (2004) proposed an algorithm for solving a two dimensional constrained cutting stock problem. In the algorithm and for depth search, hybrid approach combining hill climbing strategies and dynamic programming were employed. Cui (2005) developed an algorithm that utilizes the knapsack algorithm and an implicit enumeration technique. The algorithm was applied to real cutting stock data of the manufacture electric generators. Khalifa et al. (2006) built a one dimensional cutting stock problem using genetic algorithm. The objective was to reduce the amount of waste generated in constructing steel bars.

Saad et al. (2007) addressed the problem of scrap generated from cutting cylindrical logs produced by an aluminum extrusion company. In this regard, a multiobjective cutting stock problem was constructed. A solution procedure was developed considering the scrap generated as a fuzzy parameter. Chen (2008) presented a recursive heuristic algorithm for the constrained twodimensional stock cutting problems. The algorithm was tested and the computational results produced good solutions in short computing time for problems of different scales. Alves et al. (2009) used several constrained and non constrained integer programming using column generation. lower bounds for the different minimzation patterns were derived. Using actual data, the outcome of these models showed improvement of the lower bounds. Hajeeh (2010) addressed the problem of waste generated in an aluminum fabrication industry. A heuristic was propped for optimizing the cutting of aluminum profiles. The heuristic produced less scrap when compared to the existing procedure used in the company. Macedo et al. (2010) proposed an integer linear programming model to solve the twodimensional stock cutting problem with guillotine constraint. A computer software was used to examine the behavior of the models with data from a wood industry. The lower bound of the model was found to be superior to those of other methods.

Kasimbeyli *et al.* (2011) proposed a linear integer programing with two contarcting objectives for a onedimensional cutting problem. A special heuristic algorithm was used to find the optimim cutting pattern. Berberler et al. (2011) developed a dynamic programming algorithm to address the one-dimensional stock cutting problem. The results obtained from this algorithm was compared to others and results showed its efficeincy and superiority. Cui and Huang (2012) proposed a heuristic to address constrained T-shaped patterns with the objective of maximizing the pattern value and meeting demand. The computation of 58 benchmark instances showed that the algorithm is superior to the two-stage patterns approaches. De Valle et al. (2012) developed an algorithm based on non-fit polygen to examine the two dimensional cutting/packing problem. The algorithm aslo solved problems with items of irregular shapes. Mobasher and Ekici (2013) developed a mixed integer linear and used the column generation method to the study a cutting stock problem with set up cost. The main objective was to find a cutting pattren at minimum production cost.

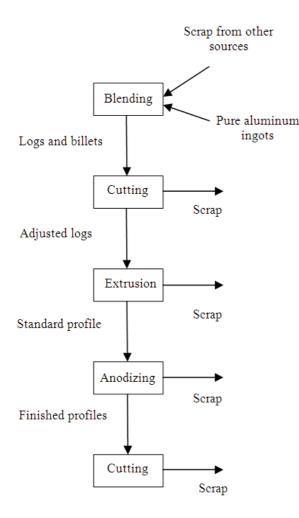
In the current research work, the extrusion process in a specific industry in Kuwait is thoroughly studied with objective of finding ways for reducing the large amount of scrap generated. The article organized as follows: it start by describing the aluminum profile production process, the amount of scrap generated in the chosen industry from using the existing cutting techniques. Next, the structure of the developed optimization models is provided. For illustration, an example is presented to compare the amount of scarp generated using the existing conventional cutting patterns and the size of scrap generate from the proposed optimization model. Results and discussion section comes next, the article ends with concluding remarks.

# 2. MATERIALS AND METHODS

## 2.1. Aluminum Profile Production

Aluminum profile production (extrusion) process passes through several stages starting with castings where logs are produced; the logs are next cut into standard billets and are put into extrusion machine to manufacture profiles of different shapes and sizes. The extruded aluminum profiles are placed in the aging furnace in order to increase their durability and strength. Next, the profiles are polished thoroughly and depending on request are either sent for paining, or anodizing before shipping to the customer. In **Fig. 1**, the detailed process is presented for a specific extrusion company in Kuwait.





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Fig. 1. Aluminum profile production process

The type of billets and logs used in the extrusion process in the company along with their lengths and weight is shown in **Table 1**. The monthly weight and percentage of scrap generated during a specific year in the aluminum extrusion by the same company is as given in **Table 2**.

#### 2.2. Optimization Models

An efficient model for minimizing scrap in the profile production process is based on developing cutting patterns, where each cutting pattern uses several billets and the total weight cut is less than the total weight of the log used. All possible patterns are investigated and the weight of the scrap generated is calculated. Two models are presented in this research work, details are given below.

 
 Table 1. Lengths and weights of standard billets and logs used by the aluminum extrusion company

Billet	Length	Weight	Log	Length	Weight
Type	(m)	(kg)	Type	(m)	(kg)
1	0.48	32.16	1	2.59	173.4
2	0.52	34.25	2	2.64	176.8
3	0.56	37.55	3	2.69	180.2
4	0.58	38.78	4	2.74	183.6
5	0.61	40.55	5	2.79	187.0
6	0.65	43.55	6	2.84	190.4

 Table 2. Weight and percentage of monthly total production scrapped in the extrusion stage

	Total	Scrap	Percentage
	production	weight	scrap
Month	(kg)	(kg)	(%)
January	249255	64722	26
February	280077	75978	27
March	295545	77954	27
April	343386	91102	27
May	245831	65408	27
June	321935	68588	21
July	268012	61195	23
August	219373	52435	24
September	317112	78005	25
October	278610	63709	23
November	347567	83844	24
December	385956	88688	23
Mean	296055	72636	25
Standard deviation	48643	11328	

## Model I

Two models are presented in this research work. The first is shown in (1) which is based on different billets:

Minimize  $Z = W_1 + W_2 + W_3$ 

Subject to Equation 1:

$$\begin{aligned} \nu_{1}\xi_{1} + \nu_{2}\xi_{1} + \nu_{3}\xi_{3} + \nu_{4}\xi_{4} + \nu_{5}\xi_{5} + \nu_{6}\xi_{1} &= W_{1} \\ \omega_{1}\xi_{1} + \omega_{2}\xi_{2} + \omega_{3}\xi_{3} + \omega_{4}\xi_{4} + \omega_{5}\xi_{5} + \omega_{6}\xi_{6} - W_{L}y &= W_{2} \\ \omega_{1}\xi_{1} + \omega_{2}\xi_{2} + \omega_{3}\xi_{3} + \omega_{4}\xi_{4} + \omega_{5}\xi_{5} + \omega_{6}\xi_{6} - W_{F}y &= W_{3} \\ y,\xi_{1},\xi_{2},\xi_{3},\xi_{4},\xi_{5},\xi_{5},\xi_{6} &\geq 0 \text{ and integers} \end{aligned}$$
(1)



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Where:

- Z = Total weight of the scarp in kg
- $W_1$  = Weight of scrap (kg) generated per 1.5 inch log (0.0381 m) of each billet used to extrude the desired the desired profile
- $W_2$  = Weight of scrap (kg) generated from log cutting
- $W_3$  = Weight of scrap (kg) generated from producing longer profiles than demanded
- $v_j$  = Weight (kg) of producing type j billet, j = 1,...,6
- $\omega_j$  = Weight (kg) of log used to produced the j<sup>th</sup> billet, j = 1,...,6
- $\xi_i$  = Number of type j billet used
- y = Number of logs used
- $W_L$  = Weight (kg) of each log
- $W_F$  = Total weight (kg) of profiles demanded

#### Model II

Although the above problem provides a good solution to the problem, however it has one disadvantage in that, where more than on log is needed, the model considers all logs to be one long log. A superior and more efficient model is based on cutting patterns as given in (2). Equation (3) represents the nenegtivity constraint.

Minimize 
$$Z = \sum_{j=1}^{J} S_j \rho_j$$

Subject to Equation 2:

$$\sum_{j=1}^{6} Q_{j} \rho_{j} \ge Q_{ML}$$
<sup>(2)</sup>

 $\rho_i$  non negative integers (3)

Where:

- $\rho_i$  = Number of cutting pattern j, j = 1, ...,6
- $S_j$  = Amount of scrap generated from cutting pattern j, j = 1, ...,6

 $Q_j$  = Weight of cutting pattern j (kg), j = 1, ...,6

 $Q_{ML}$  = Total weight of profiles demanded

The different cutting billet patterns of 1.5 m log used in the mathematical model are given in **Table 3** 

along with the total length of the different billet combinations (patterns) and the amount of scrap generated in meters.

#### 2.3. Example

As an example, the proposed model II has been used on die number 158B/ 7 for a demand 441 kg. Detailed mathematical programming formulation is as follows noting that Z and Q represent the amount of scrap generated from using two logs:

Minimize (Z + Q):

Subject to:

 $7.04X_1 + 5.04 \ X_2 + 34.04 \ X_3 + 33.04X_4 + 31.04 \ X_5$  $+28.04 X_{6} +35.07 X_{7} +32.04 X_{8} +31.04 X_{9} +29.04 X_{10}$  $+26.04X_{11} + 29.04 \ X_{12} + 28.04 \ X_{13} \ + 26.04 \ X_{14} + 23.04$  $X_{15}$  +27.04  $X_{16}$  + 25.04 $X_{17}$  +22.04  $X_{18}$  +23.04  $X_{19}$  $+20.04 X_{20} + 17.04 X_{21} - Z = 0$ 96  $X_1$  +98  $X_2$  +69  $X_3$  +604 $X_4$  +72  $X_5$  + 72  $X_6$  +68 $X_7$  $+71X_8 + 72X_9 + 74 X_{10} + 77X_{11} + 74 X_{12} + 75 X_{13} + 77 X_{14}$  $+ \hspace{0.5cm} 80 \hspace{0.1cm} X_{15} \hspace{0.1cm} + \hspace{-0.1cm} 76 \hspace{0.1cm} X_{16} \hspace{0.1cm} + \hspace{-0.1cm} 784 \hspace{0.1cm} X_{17} \hspace{0.1cm} + \hspace{-0.1cm} 81 \hspace{0.1cm} X_{18} \hspace{0.1cm} + \hspace{-0.1cm} 80 \hspace{0.1cm} X_{19} \hspace{0.1cm} + \hspace{-0.1cm} 83 \hspace{0.1cm} X_{20}$  $+ 86 X_{21} - Q = 0$ END GIN  $X_1$  $X_2$ GIN GIN  $X_3$ GIN  $X_4$ GIN  $X_5$ GIN  $X_6$ GIN  $X_7$ 

GIN	$X_8$
GIN	$X_9$
GIN	$X_{10}$
GIN	$X_{11}$
GIN	X <sub>12</sub>
GIN	X <sub>13</sub>
GIN	$X_{14}$
GIN	$X_{15}$
GIN	$X_{16}$
GIN	$X_{17}$
GIN	$X_{18}$
GIN	X19
GIN	$X_{20}$
GIN	$X_{21}$

Objective Function Value: 74.2.

	Standard billets							
Cutting patterns							Total	Total
(γ <sub>j</sub> )	0.48	0.52	0.56	0.58	0.61	0.65	length (m)	scrap (m)
$\gamma_1$	3						1.44	0.174
$\gamma_2$	2	1					1.48	0.134
γ <sub>3</sub>	1		1				1.04	0.536
$\gamma_4$	1			1			1.06	0.516
γ5	1				1		1.09	0.486
$\gamma_6$	1					1	1.13	0.446
$\gamma_7$	1	2					1.04	0.536
$\gamma_8$		1	1				1.08	0.486
γ <sub>9</sub>		1		1			1.10	0.476
$\gamma_{10}$		1			1		1.13	0.446
$\gamma_{11}$		1				1	1.17	0.406
γ <sub>12</sub>			2				1.12	0.456
γ <sub>13</sub>			1	1			1.14	0.436
$\gamma_{14}$			1		1		1.17	0.406
γ15			1			1	1.21	0.366
$\gamma_{16}$				2			1.16	0.416
$\gamma_{17}$				1	1		1.19	0.346
$\gamma_{18}$				1		1	1.23	0.356
γ19					2		1.22	0.316
γ <sub>20</sub>					1	1	1.26	0.134
<u>γ</u> 21						2	1.30	0.276

Table 3. Scrap generated from different cutting patterns ( $\gamma_i$ ) of 1.5 m Log (100.5 kg)

Total Scrap (kg.) = 100-total weight +2.5( number of billets used)

			Total scrap generated (cr	n)	
	Die	Weight	Conventional	Optimization	
No	number	demanded (kg)	method	model	
1	158 B/7	441.00	275.350	74.215	
2	861 A/3	450.00	123.400	65.200	
3	861 A/3	675.00	204.350	42.280	
4	113 D/12	396.00	226.960	119.120	
5	7613	153.00	63.720	52.400	
6	2211	165.00	65.510	41.080	
7	1096	259.00	185.430	49.920	
8	1037A/6	2574.00	311.360	208.080	
9	863A/3	888.00	76.830	64.320	
10	126A/B	244.80	328.120	205.900	
11	463	130.05	196.400	146.560	
	Total Scrap	2034.99	1145.105		

## **3. RESULTS**

**Table 3** provides the amount of scrap generated by different cutting patterns. The amount of scarp generated is less than one meter. As shown pattern 3 produces the 0.536 meters, while patterns 2 and 20 produce around

0.134 meters which is the least. **Table 4** presents the amount of scrap generated resulted for ten Dies along with the total weight demand. The amount of scrap generated using the conventional method used in the company is compared to that using the optimization model. The maximum scrap is generated by Die 1037A/6



with an amount of 208.08 kg while the least is produce by 2211 with around 41 kg.

## 4. DISCUSSION

When comparing the total amount of scrap generated for the different Dies as shown in **Table 4**, it is found that it is around 56% on average. The total scrap generated produced using the conventional procedure is around 2035 cm whereas it is around 1145 using the suggested optimization model.

# **5. CONCLUSION**

The aluminum extrusion process produces a sizable amount of scarp; this mainly attributed to the techniques used and the lack of modern scientific experience of the staff. In order to reduce the large amount of scrap, a thorough evaluation of the exiting cutting method should be carrying out. In addition to reduce scrap, an effective an optimal cutting method will contribute to efficient uses of time and other resources. For example the use of the mathematical developed, the reduction of the amount of scrap generated ranged from 21-82% and this constitutes large saving.

Efficient scientific approaches and tools should be used in the different processes within industries; optimization techniques are one of the strong tools that produce good results. It is recommended that the extrusion company should substitute the existing methods with more cost effective approaches which are based on sound scientific ones.

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